Minicourse Outline

● Lecture 1
  Basic Cilk programming: Cilk keywords, performance measures, scheduling.

● Lecture 2
  Analysis of Cilk algorithms: matrix multiplication, sorting, tableau construction.

● Laboratory
  Programming matrix multiplication in Cilk — Dr. Bradley C. Kuszmaul

● Lecture 3
  Advanced Cilk programming: inlets, abort, speculation, data synchronization, & more.
LECTURE 3

• Inlets
• Abort
• Speculative Computing
• Data Synchronization
• Under the Covers
• JCilk
• Conclusion
Operating on Returned Values

Programmers may sometimes wish to incorporate a value returned from a spawned child into the parent frame by means other than a simple variable assignment.

Example:

```
x += spawn foo(a,b,c);
```

Cilk achieves this functionality using an internal function, called an *inlet*, which is executed as a secondary thread on the parent frame when the child returns.
Semantics of Inlets

- The **inlet** keyword defines a **void** internal function to be an inlet.
- In the current implementation of Cilk, the inlet definition may not contain a **spawn**, and only the first argument of the inlet may be spawned at the call site.
Semantics of Inlets

```c
int max, ix = -1;
inlet void update ( int val, int index ) {
    if (idx == -1 || val > max) {
        ix = index; max = val;
    }
}
for (i=0; i<1000000; i++) {
    update ( spawn foo(i), i );
}
sync; /* ix now indexes the largest foo(i) */
```

1. The non-`spawn` args to `update()` are evaluated.
2. The Cilk procedure `foo(i)` is spawned.
3. Control passes to the next statement.
4. When `foo(i)` returns, `update()` is invoked.
Semantics of Inlets

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inlet void update ( int val, int index ) {
    if (idx == -1 || val > max) {
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    }
}
for (i=0; i<1000000; i++) {
    update ( spawn foo(i), i );
}sync; /* ix now indexes the largest foo(i) */
```

Cilk provides *implicit atomicity* among the threads belonging to the same frame, and thus no locking is necessary to avoid data races.
For assignment operators, the Cilk compiler automatically generates an *implicit inlet* to perform the update.
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Computing a Product

\[ p = \prod_{i=0}^{n} A_i \]

```c
int product(int *A, int n) {
    int i, p=1;
    for (i=0; i<n; i++) {
        p *= A[i];
    }
    return p;
}
```

**Optimization:** Quit early if the partial product ever becomes 0.
Computing a Product

\[ p = \prod_{i=0}^{n} A_i \]

int product(int *A, int n) {
    int i, p=1;
    for (i=0; i<n; i++) {
        p *= A[i];
        if (p == 0) break;
    }
    return p;
}

Optimization: Quit early if the partial product ever becomes 0.
Computing a Product in Parallel

\[ p = \prod_{i=0}^{n} A_i \]

```cilk
int prod(int *A, int n) {
    int p = 1;
    if (n == 1) {
        return A[0];
    } else {
        p *= spawn product(A, n/2);
        p *= spawn product(A+n/2, n-n/2);
        sync;
        return p;
    }
}
```

How do we quit early if we discover a zero?
Cilk’s Abort Feature

```cilk
int product(int *A, int n) {
    int p = 1;
    inlet void mult(int x) {
        p *= x;
        return;
    }

    if (n == 1) {
        return A[0];
    } else {
        mult( spawn product(A, n/2) );
        mult( spawn product(A+n/2, n-n/2) );
        sync;
        return p;
    }
}
```

1. Recode the implicit inlet to make it explicit.
Cilk’s Abort Feature

2. Check for 0 within the inlet.

```cilk
int product(int *A, int n) {
    int p = 1;
    inlet void mult(int x) {
        p *= x;
        return;
    }

    if (n == 1) {
        return A[0];
    } else {
        mult( spawn product(A, n/2) );
        mult( spawn product(A+n/2, n-n/2) );
        sync;
        return p;
    }
}
```
Cilk’s Abort Feature

2. Check for 0 within the inlet.
Cilk’s Abort Feature

cilk int product(int *A, int n) {
    int p = 1;
    inlet void mult(int x) {
        p *= x;
        if (p == 0) {
            abort; /* Aborts existing children, */
            /* but not future ones. */
            return;
        }
    }
    if (n == 1) {
        return A[0];
    } else {
        mult( spawn product(A, n/2) );
        mult( spawn product(A+n/2, n-n/2) );
        sync;
        return p;
    }
}
Cilk’s Abort Feature

Cilk’s Abort Feature

Implicit atomicity eases reasoning about races.
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Min-Max Search

- Two players: MAX □ and MIN ●.
- The game tree represents all moves from the current position within a given search depth.
- At leaves, apply a static evaluation function.
- MAX chooses the maximum score among its children.
- MIN chooses the minimum score among its children.
Alpha-Beta Pruning

**IDEA:** If MAX ■ discovers a move so good that MIN ● would never allow that position, MAX’s other children need not be searched — *beta cutoff.*
**Alpha-Beta Pruning**

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**Alpha-Beta Pruning**

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Alpha-Beta Pruning

**Idea:** If MAX ■ discovers a move so good that MIN ● would never allow that position, MAX’s other children need not be searched — *beta cutoff*. 

![Alpha-Beta Pruning Diagram](image-url)
IDEA: If MAX $\blacksquare$ discovers a move so good that MIN $\bullet$ would never allow that position, MAX’s other children need not be searched — *beta cutoff*. 
Alpha-Beta Pruning

IDEA: If MAX discovers a move so good that MIN would never allow that position, MAX’s other children need not be searched — beta cutoff.

Unfortunately, this heuristic is inherently serial.
**Parallel Min-Max Search**

**Observation:** In a best-ordered tree, the degree of every internal node is either 1 or *maximal*.

**Idea:** [Feldman-Mysliwietz-Monien 91] If the first child fails to generate a cutoff, *speculate* that the remaining children can be searched in parallel without wasting any work: “young brothers wait.”
Parallel Alpha-Beta (I)

cilk int search(position *prev, int move, int depth) {
    position cur;          /* Current position */
    int bestscore = -INF;  /* Best score so far */
    int num_moves;         /* Number of children */
    int mv;               /* Index of child */
    int sc;               /* Child’s score */
    int cutoff = FALSE;   /* Have we seen a cutoff? */

    ● View from MAX’s perspective; MIN’s viewpoint can be obtained by negating scores — \textit{negamax}.
    ● The node generates its current position from its parent’s position \texttt{prev}.
    ● The \texttt{alpha} and \texttt{beta} limits and the move list are fields of the \texttt{position} data structure.

#Cilk keywords used so far
Parallel Alpha-Beta (II)

```c
inlet void get_score(int child_sc) {
    child_sc = -child_sc;  /* Negamax */

    if (child_sc > bestscore) {
        bestscore = child_sc;
        if (child_sc > cur.alpha) {
            cur.alpha = child_sc;
            if (child_sc >= cur.beta) { /* Beta cutoff */
                cutoff = TRUE;  /* No need to search more */
                abort;  /* Terminate other children */
            }
        }
    }
}
```
/* Create current position and set up for search */

make_move(prev, move, &cur);

sc = eval(&cur);    /* Static evaluation */
if ( abs(sc) >= MATE || depth <= 0 ) {
    /* Leaf node */
    return (sc);
}

cur.alpha = -prev->beta;    /* Negamax */
cur.beta = -prev->alpha;

/* Generate moves, hopefully in best-first order*/
num_moves = gen_moves(&cur);
Parallel Alpha-Beta (IV)

/* Search the moves */

for (mv=0; !cutoff && mv<num_moves; mv++) {
    get_score( spawn search(&cur, mv, depth-1) );
    if (mv==0) sync; /* Young brothers wait */
}

sync;
return (bestscore);

- Only 6 Cilk keywords need be embedded in the C program to parallelize it.
- In fact, the program can be parallelized using only 5 keywords at the expense of minimal obfuscation.
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Mutual Exclusion

Cilk’s solution to mutual exclusion is no better than anybody else’s.

Cilk provides a library of spin locks declared with \texttt{Cilk\_lockvar}.

• To avoid deadlock with the Cilk scheduler, a lock should only be held within a Cilk thread.
• \textit{I.e.,} \texttt{spawn} and \texttt{sync} should not be executed while a lock is held.

Fortunately, Cilk’s control parallelism often mitigates the need for extensive locking.
Cilk’s Memory Model

Programmers may also synchronize through memory using lock-free protocols, although Cilk is agnostic on consistency model.

• If a program contains no data races, Cilk effectively supports sequential consistency.
• If a program contains data races, Cilk’s behavior depends on the consistency model of the underlying hardware.

To aid portability, the `Cilk_fence()` function implements a memory barrier on machines with weak memory models.
Cilk’s *Nondeterminator* debugging tool provably guarantees to detect and localize data-race bugs.

A *data race* occurs whenever two logically parallel threads, holding no locks in common, access the same location and one of the threads modifies the location.
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The \texttt{cilkc} compiler encapsulates the process.

\texttt{cilkc2c} translates straight C code into identical C postsource.

\texttt{gcc} translates \texttt{cilkc2c} outputs to object code.

\texttt{ld} links \texttt{gcc} object code into binary.

\texttt{Cilk RTS} is the runtime system for \texttt{Cilk} programs.
The `cilk2c` translator generates two “clones” of each Cilk procedure:

- **fast clone**—serial, common-case code.
- **slow clone**—code with parallel bookkeeping.

<table>
<thead>
<tr>
<th>SLOW</th>
<th>FAST</th>
<th>FAST</th>
<th>FAST</th>
<th>FAST</th>
</tr>
</thead>
</table>

- The **fast clone** is always spawned, saving live variables on Cilk’s work deque (shadow stack).
- The **slow clone** is resumed if a thread is stolen, restoring variables from the shadow stack.
- A check is made whenever a procedure returns to see if the resuming parent has been stolen.
Compiling `spawn` — Fast Clone

**Cilk source**

\[ x = \text{spawn} \ fib(n-1); \]

**C post-source**

```
frame->entry = 1;
frame->n = n;
push(frame);

x = fib(n-1);
if (pop() == FAILURE) {
    frame->x = x;
    frame->join--;
    \{ clean up & return to scheduler \}
}
```
Compiling `sync` — Fast Clone

No synchronization overhead in the fast clone!
Compiling the Slow Clone

```c
void fib_slow(fib_frame *frame) {
    int n,x,y;
    switch (frame->entry) {
        case 1: goto L1;
        case 2: goto L2;
        case 3: goto L3;
    }
    frame->entry = 1;
    frame->n = n;
    push(frame);
    x = fib(n-1);
    if (pop()==FAILURE) {
        frame->x = x;
        frame->join--;
        // clean up &
        // return to scheduler
    }
    if (0) {
        L1:;
        n = frame->n;
    }
}
```

- **entry**: restore program counter
- **join**: same as fast clone
- **n**: restore local variables
- **x**: if resuming
- **y**: continue

Cilk deque
Breakdown of Work Overhead (circa 1997)

Benchmark: \texttt{fib} on one processor.
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The JCilk System

- Jgo = Java + goto.
- The Jgo compiler was built by modifying gcj to accept goto statements so that a continuation mechanism for JCilk could be implemented.
JCilk Keywords

cilk
spawn sync SYNCHED
inlet abort

Same as Cilk, except that cilk can also modify try.

Eliminated!

JCilk leverages Java’s exception mechanism to render two Cilk keywords unnecessary.
Exception Handling in Java

“During the process of throwing an exception, the Java virtual machine **abruptly completes**, one by one, any expressions, statements, method and constructor invocations, initializers, and field initialization expressions that have begun but not completed execution in the current thread. This process continues until a handler is found that indicates that it handles that particular exception by naming the class of the exception or a superclass of the class of the exception.”

Exception Handling in JCilk

```java
private cilk void foo() throws IOException {
    spawn A();
    cilk try {
        spawn B();
        cilk try {
            spawn C();
        } catch (ArithmeticException e) {
            doSomething();
        }
    } catch (RuntimeException e) {
        doSomethingElse();
    }
    spawn D();
    doYetSomethingElse();
sync;
}
```
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        doYetSomethingElse();
        sync;
    }

An exception causes all subcomputations dynamically enclosed by the catching clause to abort!
private cilk void foo() throws IOException {
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        doSomethingElse();
    }
    spawn D();
    doYetSomethingElse();
    sync;
}
JCilk’s Exception Mechanism

- JCilk’s exception semantics allow programs such as alpha-beta to be coded without Cilk’s *inlet* and *abort* keywords.

- Unfortunately, Java exceptions are slow, reducing the utility of JCilk’s faithful extension.
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Future Work

Adaptive computing

• Get rid of --nproc.
• Build a job scheduler that uses parallelism feedback to balance processor resources among Cilk jobs.

Integrating Cilk with static threads

• Currently, interfacing a Cilk program to other system processes requires arcane knowledge.
• Build linguistic support into Cilk for Cilk processes that communicate.
• Develop a job scheduler that uses pipeload to allocate resources among Cilk processes.
Key Ideas

- Cilk is simple: `cilk, spawn, sync, SYNCHED, inlet, abort`
- JCilk is simpler
- Work & span
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Open-Cilk Consortium

- We are in the process of forming a consortium to manage, organize, and promote Cilk open-source technology.
- If you are interested in participating, please let us know.